

Convex Optimization

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Lecture 3:

Steepest Descent
and Gradient Descent

Reading: Boyd and Vandenberghe 9.1-9.4

Alternative reading: Bertsekas Nonlinear Programming 1.2-1.3

Optional reading on other linesearch and stepsize strategies:

Nocedal and Wright 3.1-3.2, parts of 3.3, 3.5

Unconstrained Optimization Problems

$$\min_{x \in \mathbb{R}^n} f(x)$$

$$f: \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$$

- We will assume (for now) that we are given some starting point $x^{(0)} \in \text{dom}(f)$ (i.e. with $f(x^{(0)}) < \infty$)
- Optimize given $x^{(0)}$ and access to a 1st order oracle
 $x \mapsto f(x), \nabla f(x)$

Center of Mass Method

- Requires keeping track of polyhedral with increasing number of facets— $O(nk) = O\left(n^2 \log \frac{1}{\epsilon}\right)$ memory
- Requires computing center of mass
 - Equivalent to integrating—harder than optimizing!
 - Can approximate well enough in randomized $\text{poly}(n)$ time
- Reasonable number of iterations / grad evals:
 $O\left(n \log \frac{1}{\epsilon}\right)$
...but horrible runtime

(Generic) Descent Method

Init $x^{(0)} \in \text{dom}(f)$
Iterate $x^{(k+1)} \leftarrow x^{(k)} + t^{(k)} \Delta x^{(k)}$

stepsize
 $t \in \mathbb{R}$

direction
 $\Delta x \in \mathbb{R}$

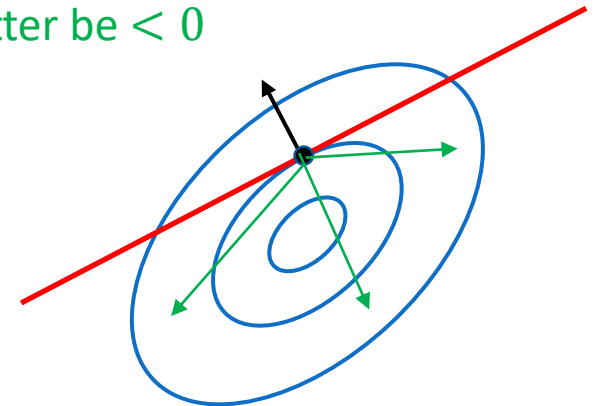
- Def: Δx is a descent direction iff $\langle \nabla f(x), \Delta x \rangle < 0$

- Recall:

$$f(x^+) = f(x + t\Delta x) \geq f(x) + t \underbrace{\langle \nabla f(x), \Delta x \rangle}_{\text{Better be } < 0}$$

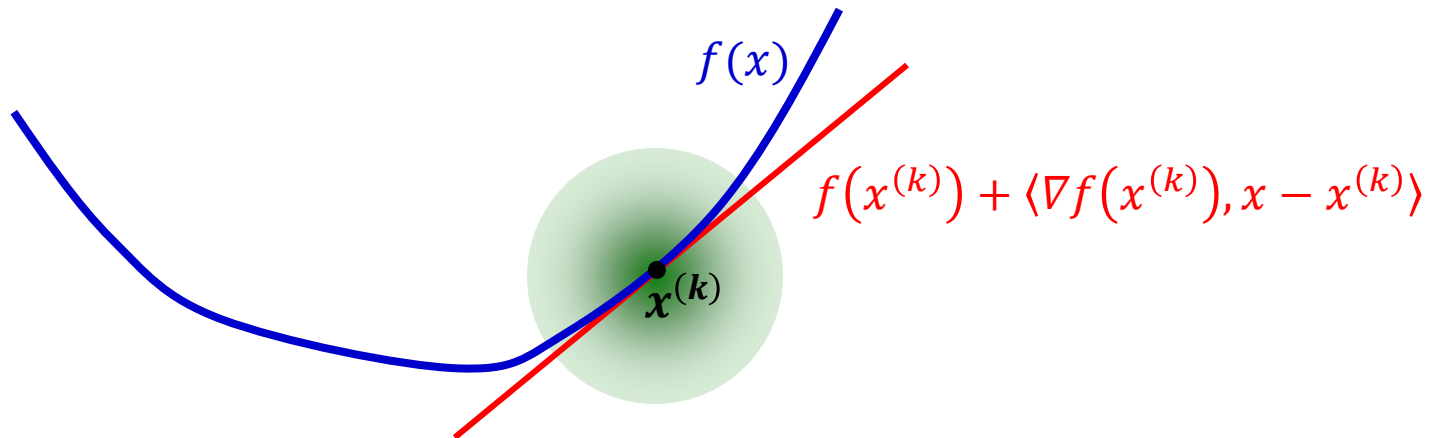
- Cutting plane view:

$$x^* \in \{x \mid \langle \nabla f(x^{(k)}), x - x^{(k)} \rangle < 0\}$$



Which Descent Direction?

$$x^{(k+1)} \leftarrow \arg \min \underbrace{f(x^{(k)}) + \langle \nabla f(x^{(k)}), x - x^{(k)} \rangle}_{\text{1st order approx of } f(x)} + \underbrace{\frac{\alpha}{2} \|x - x^{(k)}\|^2}_{\text{Only valid near } x^{(k)}}$$



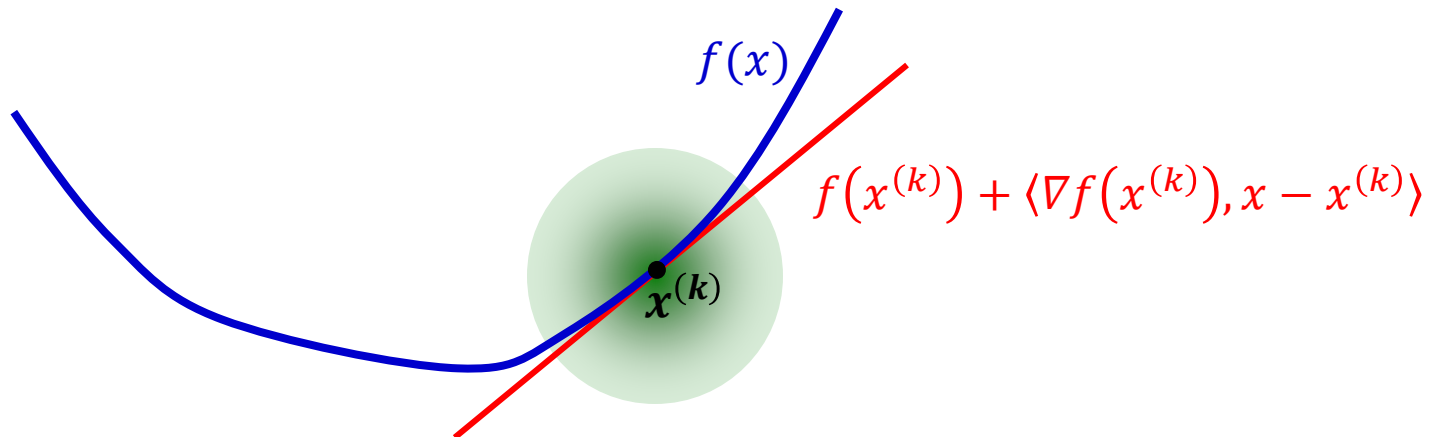
Which Descent Direction?

$$x^{(k+1)} \leftarrow \arg \min f(x^{(k)}) + \langle \nabla f(x^{(k)}), x - x^{(k)} \rangle + \frac{\alpha}{2} \|x - x^{(k)}\|^2$$
$$= \arg \min_{\|\Delta x\|^2 = 1} f(x^{(k)}) + t \langle \nabla f(x^{(k)}), \Delta x \rangle + \frac{\alpha}{2} t^2 \|\Delta x\|^2$$

$x = x^{(k)} + t\Delta x$

$$t = \frac{-\langle \nabla f(x), \Delta x \rangle}{\alpha}$$

$$\Delta x = -\arg \max_{\|v\|=1} \langle \nabla f(x^{(k)}), v \rangle$$



Direction of Steepest Descent

$$\Delta x = -\arg \max_{\|v\|=1} \langle \nabla f(x^{(k)}), v \rangle$$

- Choice of norm $\| \cdot \|$ is crucial!
- Using the Euclidean norm: $\|x\|_2 = \sqrt{\sum_i x[i]^2}$:
 $\Delta x \propto -\nabla f(x^{(k)})$ (as vectors)
- Depends on choice of basis!
- Choice of norm (eg basis for Euclidean norm) relates the primal x space to the dual space of gradients ∇f

Gradient Descent

(Steepest Descent w.r.t Euclidean Norm)

$$\Delta x = -\nabla f(x^{(k)})$$

Init $x^{(0)} \in \text{dom}(f)$

Iterate $x^{(k+1)} \leftarrow x^{(k)} - t^{(k)} \nabla f(x^{(k)})$

Reminder: We are violating here the distinction between the primal space and the dual gradient space—we are implicitly linking them by matching representations w.r.t. a chosen basis

Note: Δx is not normalized (i.e. we don't require $\|\Delta x\|_2 = 1$). This just changes the meaning of t .

How do we choose the stepsize $t^{(k)}$?

Setting the Stepsize

Option 1: Exact Linesearch

$$t^{(k)} \leftarrow \arg \min_{t \in \mathbb{R}} f(x^{(k)} + t\Delta x^{(k)})$$

- This is a convex one-dimensional problem
 ➔ can use bisection!
- But to what accuracy?
- Outer loop (updating $x^{(k)}$) and inner loop (optimizing $t^{(k)}$):
 when do we stop inner loop and iterate outer loop?

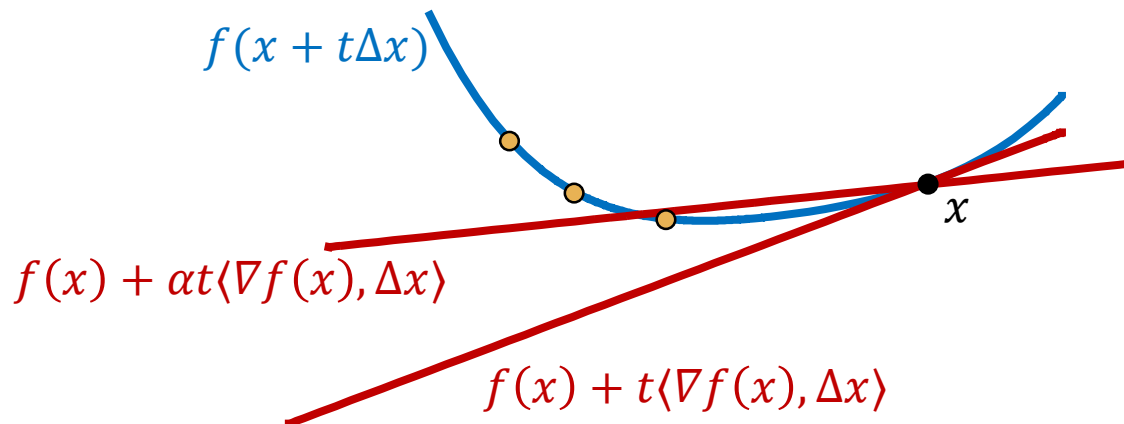
Option 2: Backtracking Linesearch (Armijo's Rule)

- Parameters: $0 < \alpha < \frac{1}{2}$ and $0 < \beta < 1$
- Input: initial point x
subgradient $\nabla f(x)$ at initial point
direction/vector Δx
evaluation oracle for $f(\cdot)$
- Output: stepsize t

Init $t \leftarrow 1$

Until $f(x + t\Delta x) < f(x) + \alpha \cdot t\langle \nabla f(x), \Delta x \rangle$

Set $t \leftarrow \beta \cdot t$



Gradient Descent

Init $x^{(0)} \in \text{dom}(f)$

Iterate $\Delta x^{(k)} = -\nabla f(x^{(k)})$

Set $t^{(k)}$ by backtracking linesearch

$$x^{(k+1)} \leftarrow x^{(k)} + t^{(k)} \Delta x^{(k)}$$

Stopping condition?

Runtime analysis? How many iterations?

Smoothness and Strong Convexity

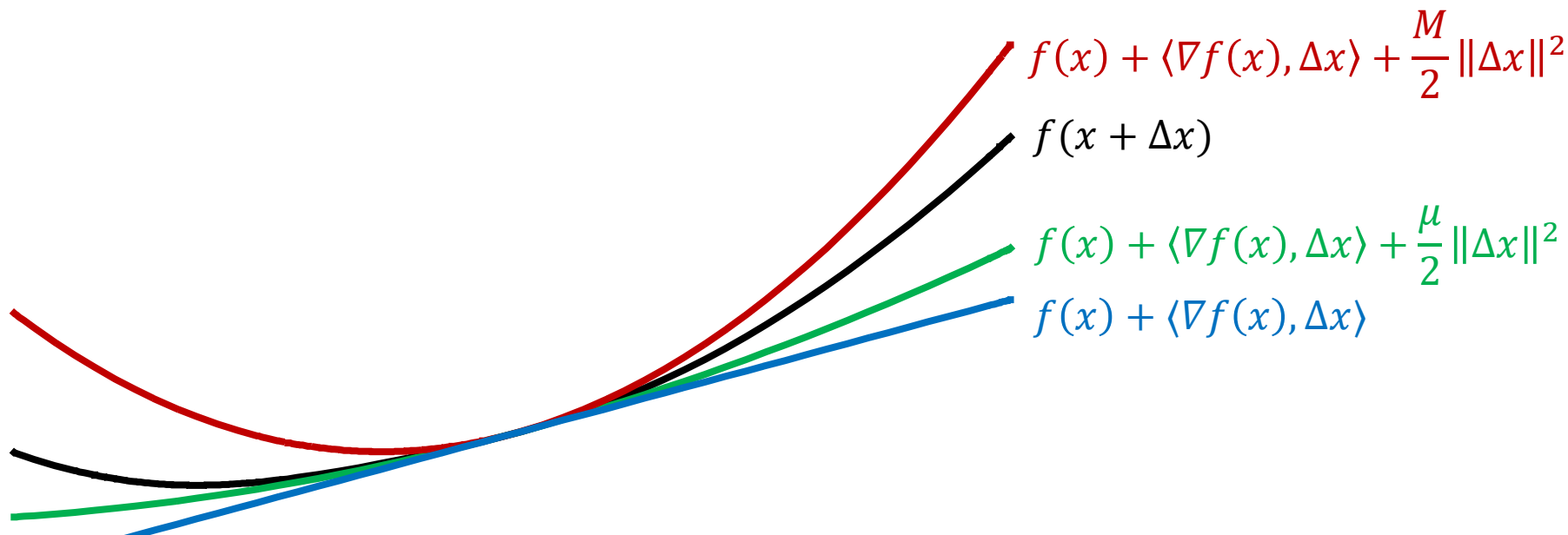
Def: f is μ -strongly convex

Def: f is M -smooth

$$f(x) + \langle \nabla f(x), \Delta x \rangle + \frac{\mu}{2} \|\Delta x\|_2^2 \leq f(x + \Delta x) \leq f(x) + \langle \nabla f(x), \Delta x \rangle + \frac{M}{2} \|\Delta x\|_2^2$$

Can be viewed as a condition on the directional 2nd derivatives

$$\mu \leq f_v''(x) = \frac{\partial^2}{\partial t^2} f(x + tv) = v^\top \nabla^2 f(x) v \leq M \quad (\text{for } \|v\|_2 = 1)$$



Smoothness and Strong Convexity

Def: f is μ -strongly convex

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$$f(x) + \langle \nabla f(x), \Delta x \rangle + \frac{\mu}{2} \|\Delta x\|_2^2 \leq f(x + \Delta x) \leq f(x) + \langle \nabla f(x), \Delta x \rangle + \frac{M}{2} \|\Delta x\|_2^2$$

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And as condition on eigenvalues of Hessian:

$$\mu \leq \lambda_{\min}(\nabla^2 f(x)), \lambda_{\max}(\nabla^2 f(x)) \leq M$$

$$\mu I \preceq \nabla^2 f(x) \preceq M I$$

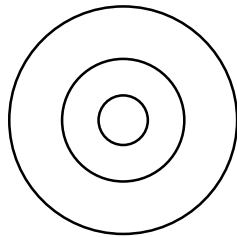
$$\kappa = \frac{M}{\mu} = \frac{\max_x \lambda_{\max}(\nabla^2 f(x))}{\min_x \lambda_{\min}(\nabla^2 f(x))}$$

Examples

$$f(x) = \frac{1}{2}x^\top Hx + b^\top x$$

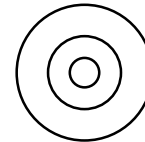
$$H = \begin{bmatrix} 1 & 0 \\ 0 & 1 \end{bmatrix}$$

$\mu = 1, M = 1$



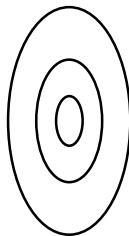
$$H = \begin{bmatrix} 5 & 0 \\ 0 & 5 \end{bmatrix}$$

$\mu = 5, M = 5$



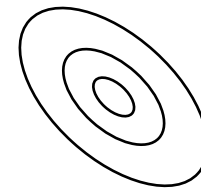
$$H = \begin{bmatrix} 5 & 0 \\ 0 & 1 \end{bmatrix}$$

$\mu = 1, M = 5$



$$H = \begin{bmatrix} 3 & 2 \\ 2 & 3 \end{bmatrix}$$

$\mu = 1, M = 5$



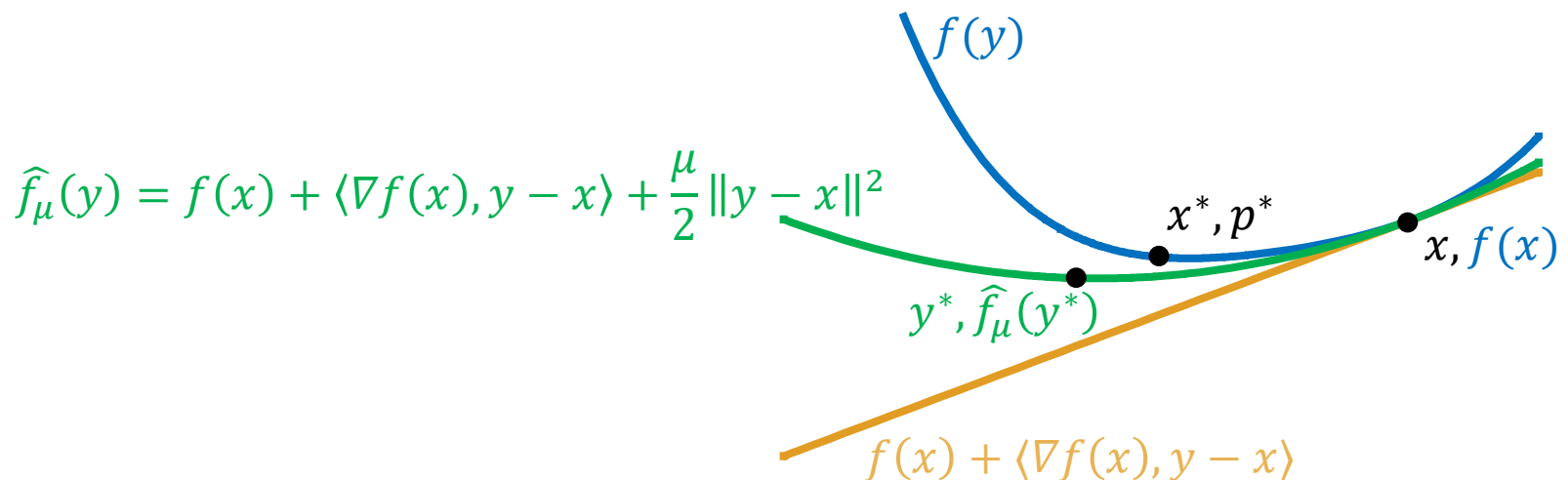
Strong Convexity and Sub-optimality

- Assume f is μ -strongly-convex (i.e. $\forall_x \mu I \preceq \nabla^2 f(x)$)

$$p^* = f(x^*)$$

$$\geq \min_y \underbrace{f(x) + \langle \nabla f(x), y - x \rangle + \frac{\mu}{2} \|y - x\|_2^2}_{\hat{f}_\mu(y)} = \underbrace{f(x) - \frac{1}{2\mu} \|\nabla f(x)\|_2^2}_{\hat{f}_\mu(y^*)}$$

$y^* = x - \frac{1}{\mu} \nabla f(x)$



Strong Convexity and Sub-optimality

- Assume f is μ -strongly-convex (i.e. $\forall_x \mu I \preceq \nabla^2 f(x)$)

$$\begin{aligned} p^* = f(x^*) &\geq f(x) + \langle \nabla f(x), x^* - x \rangle + \frac{\mu}{2} \|x^* - x\|_2^2 \\ &\geq \min_y f(x) + \langle \nabla f(x), y - x \rangle + \frac{\mu}{2} \|y - x\|_2^2 = f(x) - \frac{1}{2\mu} \|\nabla f(x)\|_2^2 \end{aligned}$$

$$\Rightarrow f(x) - p^* \leq \frac{1}{2\mu} \|\nabla f(x)\|_2^2$$

- Stopping condition to ensure ϵ -suboptimality:

$$\|\nabla f(x)\| \leq \sqrt{2\mu\epsilon}$$

Runtime Analysis in terms of $\kappa = \frac{M}{\mu}$

- Assuming strong convexity AND smoothness $\forall_x \mu I \leq \nabla^2 f(x) \leq M I$
- Using exact linesearch:

$$\begin{aligned} f(x^+) &= \min_t f(x - t \nabla f(x)) \\ &\leq \min_t f(x) + \langle \nabla f(x), -t \nabla f(x) \rangle + \frac{M}{2} \|t \nabla f(x)\|_2^2 \\ &= \min_t f(x) - t \|\nabla f(x)\|^2 + t^2 \frac{M}{2} \|\nabla f(x)\|_2^2 = f(x) - \frac{1}{2M} \|\nabla f(x)\|_2^2 \\ &\leq f(x) - \frac{2\mu}{2M} (f(x) - p^*) \end{aligned}$$

$$\Rightarrow f(x^{(k+1)}) - p^* \leq f(x^{(k)}) - p^* - \frac{1}{\kappa} (f(x^{(k)}) - p^*) = \left(1 - \frac{1}{\kappa}\right) (f(x^{(k)}) - p^*)$$

$$\Rightarrow f(x^{(k)}) - p^* \leq \left(1 - \frac{1}{\kappa}\right)^k (f(x^{(0)}) - p^*)$$

➔ Number of iterations required for $f(x^{(k)}) \leq p^* + \epsilon$:

$$k \leq \frac{1}{\log\left(\frac{\kappa}{\kappa-1}\right)} \log\left(\frac{f(x^{(0)}) - p^*}{\epsilon}\right)$$

Gradient Descent for Smooth and Strongly Convex Objectives

- Theorem: If f is μ -strongly convex and M -smooth on $S^0 = \{x \mid f(x) \leq f(x^{(0)})\}$ (i.e. $\forall_{f(x) \leq f(x^{(0)})} \mu I \preceq \nabla^2 f(x) \preceq MI$), then Gradient Descent with exact line search with at most

$$k \leq \frac{1}{-\log\left(1 - \frac{1}{\kappa}\right)} \log\left(\frac{f(x^{(0)}) - p^*}{\epsilon}\right) \approx \kappa \cdot \log\left(\frac{f(x^{(0)}) - p^*}{\epsilon}\right)$$

iterations ensures $f(x^{(k)}) \leq p^* + \epsilon$

Using Backtracking Lineasearch

Init $t \leftarrow 1$

Until $f(x + t\Delta x) \leq f(x) + \alpha \cdot t\langle \nabla f(x), \Delta x \rangle = f(x) - t\alpha \|\nabla f(x)\|_2^2$

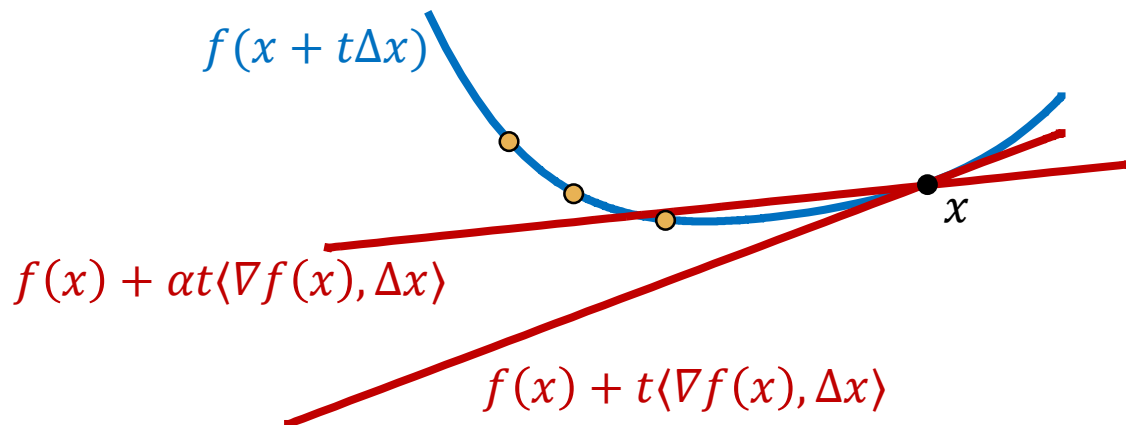
Set $t \leftarrow \beta \cdot t$

- Claim: If f is M -smooth, then any $t < \frac{1}{M}$ satisfies Armijo for any $\alpha < 1/2$

Proof: $f(x - t\nabla f(x)) \leq f(x) + \langle \nabla f(x), -t\nabla f(x) \rangle + \frac{M}{2} \|t\nabla f(x)\|_2^2$

$$= f(x) - \left(1 - \frac{M}{2}t\right)t\|\nabla f(x)\|_2^2 \leq f(x) - t\frac{1}{2}\|\nabla f(x)\|_2^2 \leq f(x) - t\alpha\|\nabla f(x)\|_2^2$$

- Conclusion: we either use $t = 1$ or $t > \beta/M$



Using Backtracking Lineasearch

Init $t \leftarrow 1$

Until $f(x + t\Delta x) \leq f(x) + \alpha \cdot t \langle \nabla f(x), \Delta x \rangle = f(x) - t\alpha \|\nabla f(x)\|_2^2$

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$$= f(x) - \left(1 - \frac{M}{2}t\right) t \|\nabla f(x)\|_2^2 \leq f(x) - t \frac{1}{2} \|\nabla f(x)\|_2^2 \leq f(x) - t\alpha \|\nabla f(x)\|_2^2$$

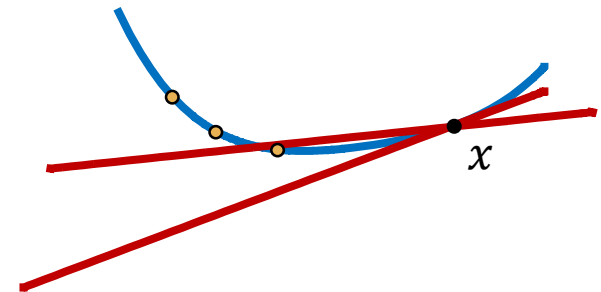
- Conclusion: we either use $t = 1$ or $t > \beta/M$

$$\rightarrow f(x - t\nabla f(x)) \leq f(x) - \min\left(1, \frac{\beta}{M}\right) \alpha \|\nabla f(x)\|_2^2$$

$$\text{\#iter } k \leq \frac{1}{-\log\left(1 - 2\alpha \min\left(\mu, \frac{\beta}{\kappa}\right)\right)} \log\left(\frac{f(x^{(0)}) - p^*}{\epsilon}\right)$$

- How many inner iterations?

Since $t < 1/M$ always OK, at most $\left\lceil \frac{\log \frac{1}{M}}{\log \beta} \right\rceil$ per outer iteration



Gradient Descent with Backtracking Linesearch

Init $x^{(0)} \in \text{dom}(f)$
Iterate $\Delta x^{(k)} = -\nabla f(x^{(k)})$
Stop if $\|\nabla f(x^{(k)})\|_2^2 \leq 2\mu\epsilon$
Set $t^{(k)}$ by backtracking linesearch with params α, β
 $x^{(k+1)} \leftarrow x^{(k)} + t^{(k)} \Delta x^{(k)}$

If f is μ -strongly-convex and M -smooth on $\{f(x) \leq f(x^{(0)})\}$, #func evals :

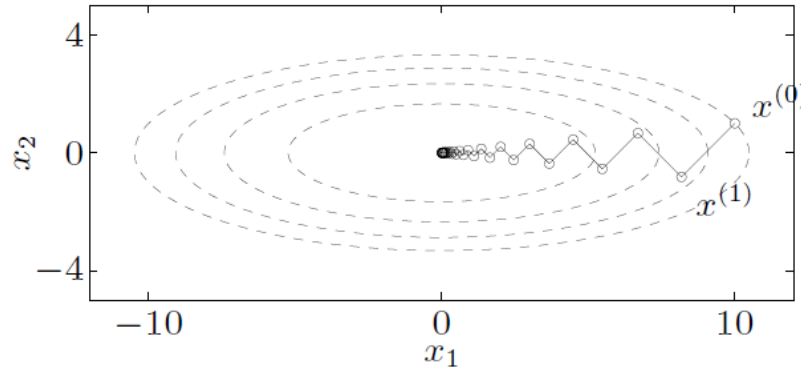
$$k = O\left(\left\lceil \frac{\log M}{\log \frac{1}{\beta}} \right\rceil \frac{1}{\alpha} \max\left(\frac{1}{M}, \frac{1}{\beta}\right) \kappa \log\left(\frac{f(x^{(0)}) - p^*}{\epsilon}\right)\right)$$

$\kappa = M/\mu$

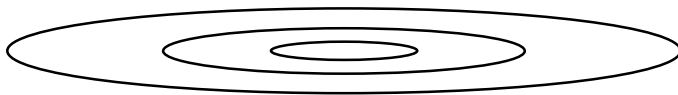
and additional runtime: $O\left(n \left\lceil \frac{\log M}{\log \frac{1}{\beta}} \right\rceil \frac{1}{\alpha} \max\left(\frac{1}{M}, \frac{1}{\beta}\right) \kappa \log\left(\frac{f(x^{(0)}) - p^*}{\epsilon}\right)\right)$

Is dependence on κ real?

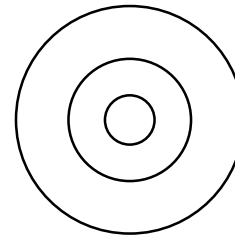
- Yes! Even for quadratic, even with exact linesearch.



- Solution:



$$f(x) = \frac{1}{2} x^\top H x$$



$$\begin{aligned} \tilde{f}(\tilde{x}) &= f(H^{-1/2} \tilde{x}) = \frac{1}{2} \tilde{x}^\top I \tilde{x} \\ \tilde{x} &= H^{1/2} x \end{aligned}$$

- Another way to view this: change of basis

Change of variables / basis

$$\tilde{f}(\tilde{x}) = f(H^{-1/2}\tilde{x}) \quad \tilde{x} \leftarrow H^{1/2}x$$

- GD direction:

$$\Delta\tilde{x} = -\nabla\tilde{f}(\tilde{x}) = -H^{-1/2}\nabla f(H^{-1/2}\tilde{x})$$

- What's the update in the original basis?

$$\tilde{x}^+ \leftarrow \tilde{x} + t\Delta\tilde{x} \quad \equiv \quad x \leftarrow x + t\Delta x$$

$$\Delta x = H^{-1/2}\Delta\tilde{x} = -H^{-1/2}H^{-1/2}\nabla f(H^{-1/2}\tilde{x}) = -H^{-1}\nabla f(x)$$

- Alternative view: steepest descent w.r.t. $\|x\|_H$

Pre-Conditioned Gradient Descent

Init $x^{(0)} \in \text{dom}(f)$

$$H = \nabla^2 f(x^{(0)})$$

Iterate $\Delta x^{(k)} = -H^{-1} \nabla f(x^{(k)})$

Set $t^{(k)}$ by backtracking linesearch with params α, β

$$x^{(k+1)} \leftarrow x^{(k)} + t^{(k)} \Delta x^{(k)}$$

